

Signal machines : localization of isolated accumulation

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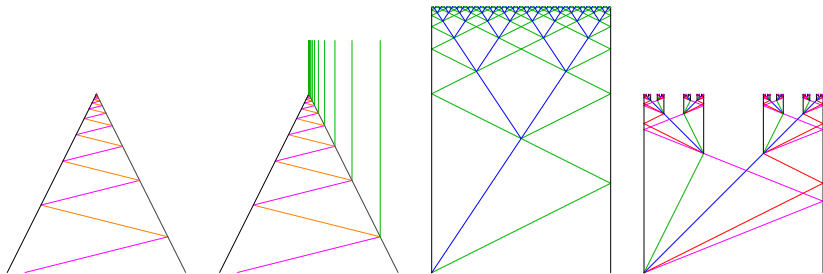


6 mars 2011 — Journées Calculabilités — Paris

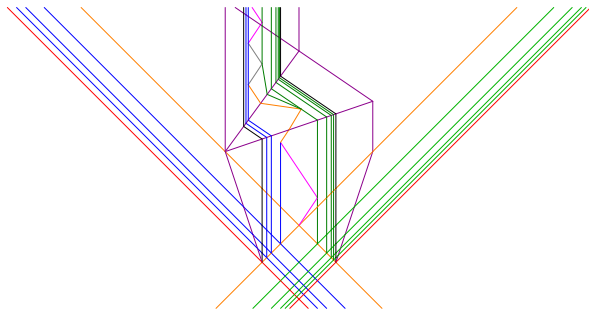
- 1 Signal machines and isolated accumulations
- 2 Necessary conditions on the coordinates of isolated accumulations
- 3 Manipulating *c.e.* and *d-c.e.* real numbers
- 4 Accumulating at *c.e.* and *d-c.e.* real numbers
- 5 Conclusion

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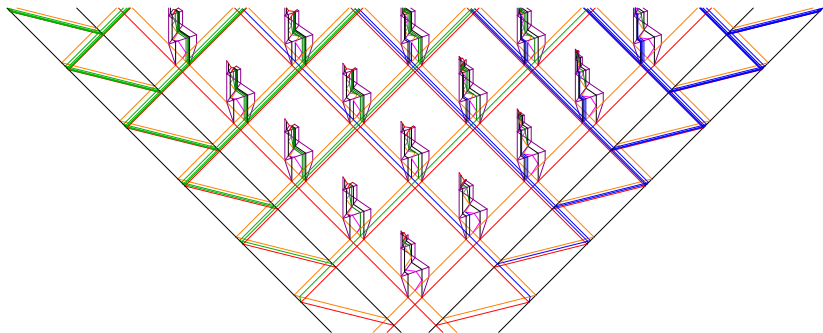
“Nice regular drawings”



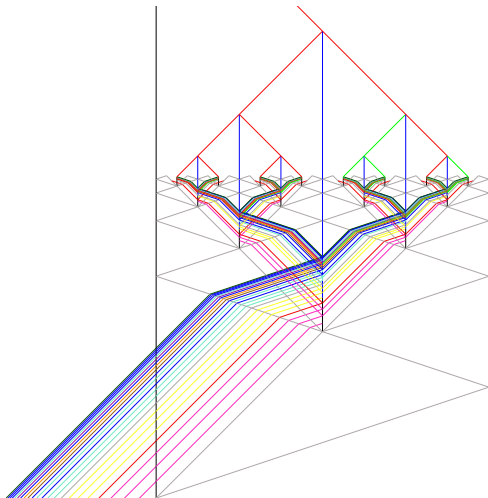
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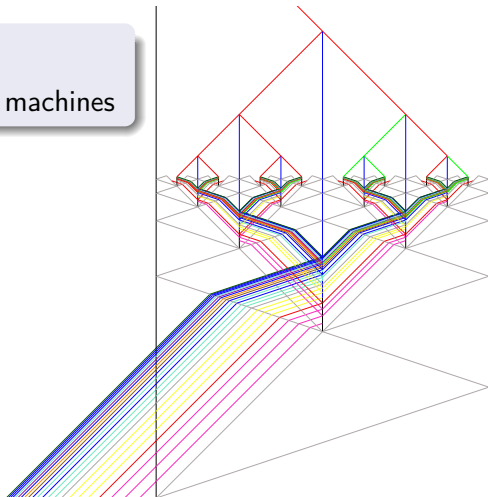
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Lines: traces of signals

Space-time diagrams of signal machines



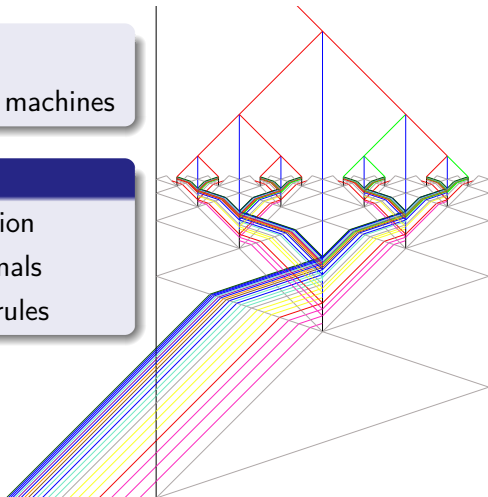
“Nice regular drawings”

Lines: traces of signals

Space-time diagrams of signal machines

Defined by

- bottom: initial configuration
- lines: signals \rightsquigarrow meta-signals
- end-points: collisions \rightsquigarrow rules



Example: find the middle

M |

M |

Meta-signals (speed)

M (0)

Collision rules

Example: find the middle



Meta-signals (speed)

M (0)
div (3)

Collision rules

Example: find the middle



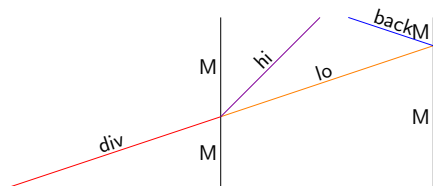
Meta-signals (speed)

M (0)
div (3)
hi (1)
lo (3)

Collision rules

$\{ \text{div}, M \} \rightarrow \{ M, \text{hi}, \text{lo} \}$

Example: find the middle



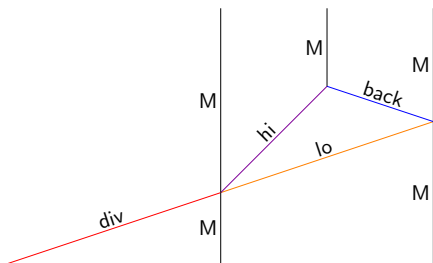
Meta-signals (speed)

M (0)
div (3)
hi (1)
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back (-3)

Collision rules

$\{ \text{div}, M \} \rightarrow \{ M, \text{hi}, \text{lo} \}$
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Example: find the middle



Meta-signals (speed)

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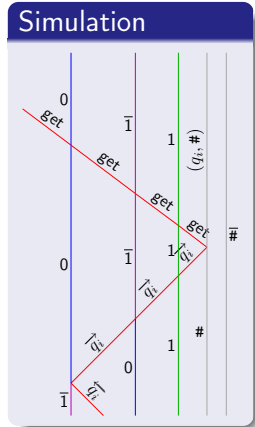
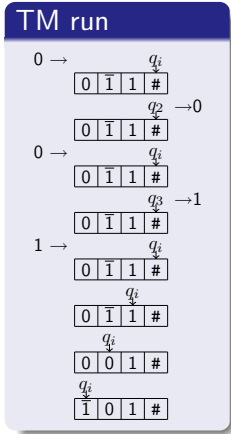
Collision rules

$\{ \text{div}, M \} \rightarrow \{ M, \text{hi}, \text{lo} \}$
 $\{ \text{lo}, M \} \rightarrow \{ \text{back}, M \}$
 $\{ \text{hi}, \text{back} \} \rightarrow \{ M \}$

Known results

Turing computations

- [Durand-Lose, 2011]



Known results

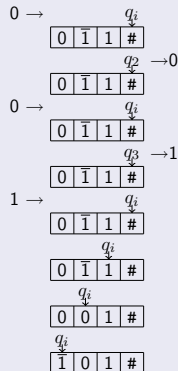
Turing computations

- [Durand-Lose, 2011]

Analog computations

- Computable analysis [Weihrauch, 2000] [Durand-Lose, 2010a]
- Blum, Shub and Smale model [Blum et al., 1989] [Durand-Lose, 2008]

TM run



Simulation



Known results

Turing computations

- [Durand-Lose, 2011]

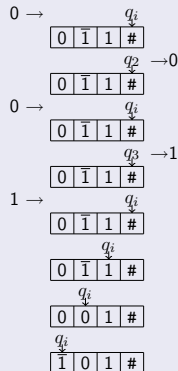
Analog computations

- Computable analysis [Weihrauch, 2000] [Durand-Lose, 2010a]
- Blum, Shub and Smale model [Blum et al., 1989] [Durand-Lose, 2008]

“Black hole” implementation

- [Durand-Lose, 2009]

TM run

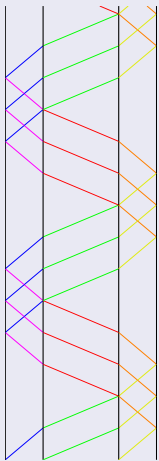


Simulation

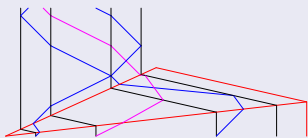


Geometric primitives: accelerating and bounding time

Normal

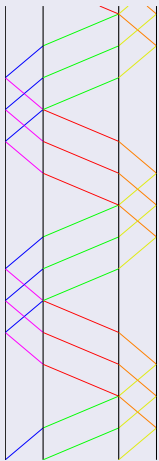


Shrunk

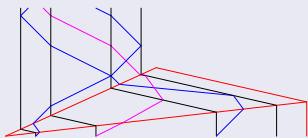


Geometric primitives: accelerating and bounding time

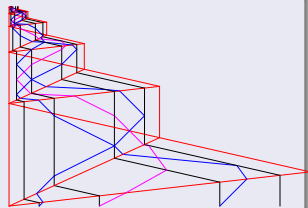
Normal



Shrunk



Iterated



Rational signal machines and isolated accumulations

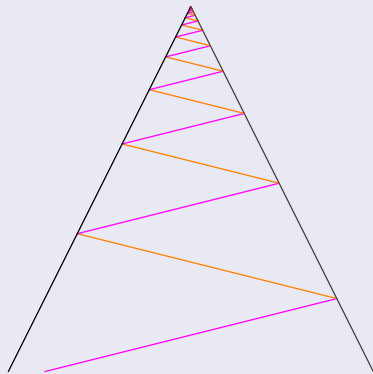
\mathbb{Q} signal machine

- all speed are in \mathbb{Q}
- all initial positions are in \mathbb{Q}
- \Rightarrow all location remains in \mathbb{Q}

Space and time location

- Easy to compute

Simplest accumulation



Rational signal machines and isolated accumulations

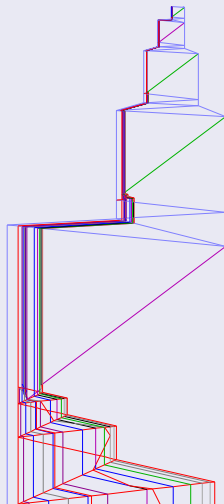
\mathbb{Q} signal machine

- all speed are in \mathbb{Q}
- all initial positions are in \mathbb{Q}
- \Rightarrow all location remains in \mathbb{Q}

Space and time location

- Easy to compute
- Not so easy to guess

Accumulation?



Rational signal machines and isolated accumulations

\mathbb{Q} signal machine

- all speed are in \mathbb{Q}
- all initial positions are in \mathbb{Q}
- \Rightarrow all location remains in \mathbb{Q}

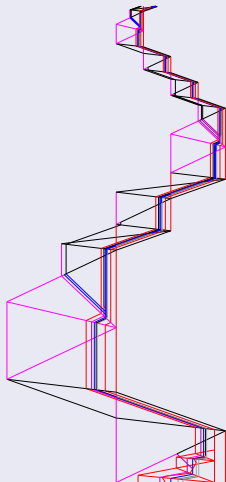
Space and time location

- Easy to compute
- Not so easy to guess

Forecasting any accumulation

Highly undecidable
(Σ_2^0 in the arithmetic hierarchy)
[Durand-Lose, 2006]

Accumulation?



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Temporal coordinate

\mathbb{Q} -signal machine

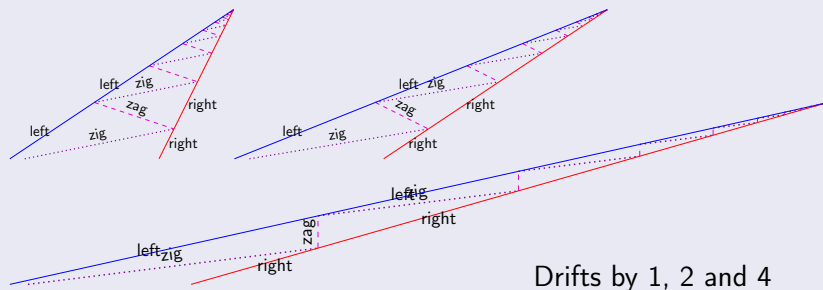
- \mathbb{Q} on computers/Turing machine
 - exact representation
 - exact operations
- exact computations by TM (and implanted in Java)

Simulation near an isolated accumulation

- on each collision, print the date
- \rightsquigarrow increasing computable sequence of rational numbers (converges iff there is an accumulation)

Spatial coordinate

Static deformation by adding a constant to each speed



With all speeds positive

- the left most coordinate is increasing (and computable) converges iff there is an accumulation
- correction by subtracting the date times the drift

c.e. real number

- limit of a convergent increasing computable sequence of rational numbers
- no bound on the convergence rate
- represents a c.e. set (of natural numbers)
- stable by positive integer multiplication but not by subtraction

d-c.e. real number

- difference of two c.e. real number
- form a field
- [Ambos-Spies et al., 2000]
these are exactly the limits of a computable sequence of rational numbers that converges *weakly effectively*, i.e.,

$$\sum_{n \in \mathbb{N}} |x_{n+1} - x_n| \text{ converges}$$

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Encoding

For *d-c.e.* real numbers

$$x = \sum_{i \in \mathbb{N}} \frac{z_i}{2^i}, z_i \in \mathbb{Z}$$

the sequence $i \rightarrow z_i$ is computable and

$$\sum_{i \in \mathbb{N}} \left| \frac{z_i}{2^i} \right| \text{ converges}$$

For *c.e.* real numbers

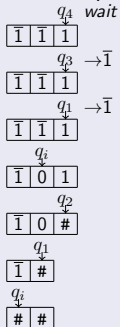
- identical but $z_i \in \mathbb{N}$
- z_i in signed unary representation

TM outputting the infinite sequence

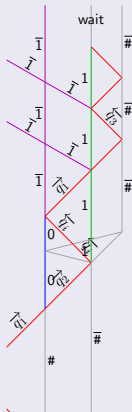
Run

wait between

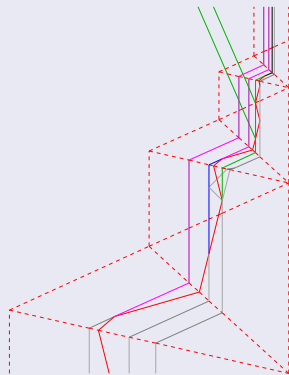
each z_i



Simulation



Shrunk to output in bounded time



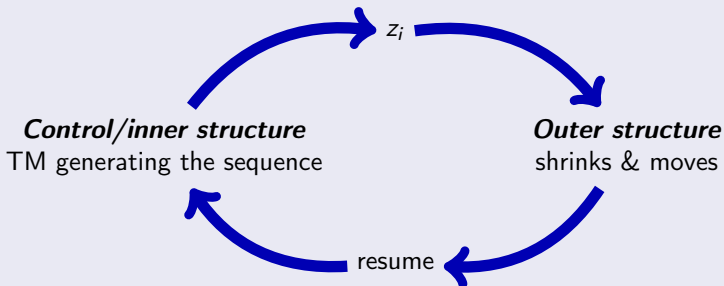
Simulation and shrinking structure stop after each value

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Two-level scheme

Control/inner structure

- Provide the data for accumulating

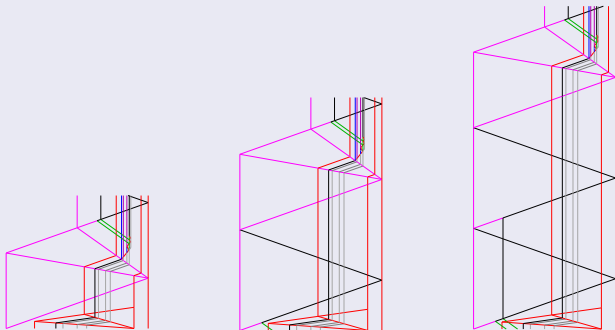


Outer structure

- shrink and move the whole structure \rightsquigarrow accumulation

Temporal coordinate

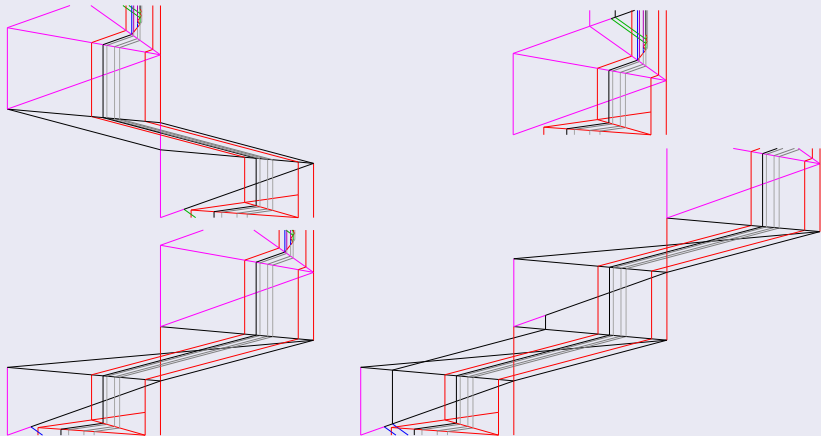
Wait the corresponding time



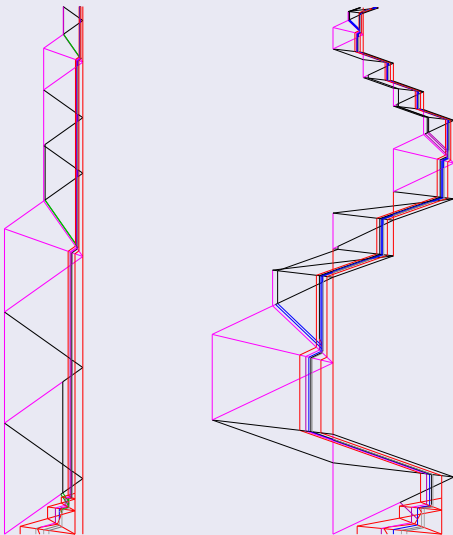
- Constant (up to scale) delay before outer structure action
- total delay is rational and should be previously subtracted

Spatial coordinate

Move left or right, more or less



Examples



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Results

- Isolated accumulations happen at *d-c.e.* spacial and *c.e.* temporal coordinates
- Accumulation at any *c.e.* temporal coordinate is possible
- Accumulation at any *d-c.e.* spacial coordinate is possible

Perspectives

- Uncorrelate space and time coordinate
it is possible for *computable* coordinates [Durand-Lose, 2010b]
- Higher order isolated accumulations
- Non isolated accumulations



Ambos-Spies, K., Weihrauch, K., and Zheng, X. (2000).

Weakly computable real numbers.

J. Complexity, 16(4):676–690.



Blum, L., Shub, M., and Smale, S. (1989).

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Applications of Models of Computations (TAMC '06)*, number 3959
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